# Digital Logic Design: a rigorous approach (C) Chapter 1: Sets and Functions

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- Naive definition of sets fails due to paradoxes (Cantor, Russel). Beginning of 20th century: axiomatization of set theory (Zermelo-Fraenkel axioms).
- Bypass based on a universal set.

The universal set is a set that contains all the possible objects.

- $\bullet$  Universal set set of all real numbers  $\mathbb R$
- Universal set set of all natural numbers  $\mathbb N$  (integers  $\geq 0$ ) numbers.

A set is a collection of objects from a universal set.

# **Specification**

We denote the set of all elements in U that satisfy property  $P$  as follows

$$
\{x \in U \mid x \text{ satisfies property } P\}.
$$

# Notation: the symbol  $\triangleq$

 $\mathbb{N}^+ \stackrel{\scriptscriptstyle\triangle}{=} \{n\in\mathbb{N} \mid n\geq 1\}$  means " $\mathbb{N}^+$  is defined be the set of all positive natural numbers". (Compare:  $=$  and  $\stackrel{\triangle}{=}$ )

- $\mathbb{Q} \stackrel{\scriptscriptstyle\triangle}{=} \{ \mathsf{x} \in \mathbb{R} \mid \mathsf{x} \text{ is a rational number} \}$
- $P \stackrel{\triangle}{=} \{x \in \mathbb{N} \mid x \text{ is a prime number}\}$

$$
\bullet\,\,\mathbb{Z}\stackrel{\triangle}{=}\{x\in\mathbb{R}\mid x\,\,\text{is a multiple of}\,\,1\}
$$

- $\mathbb{N} \stackrel{\triangle}{=} \{x \in \mathbb{Z} \mid x \geq 0\}$
- set of even integers is  $\{x \in \mathbb{Z} \mid x \text{ is a multiple of } 2\}$
- Suppose  $U \stackrel{\triangle}{=} \mathbb{N}$ .
- $\mathcal{A}\stackrel{\scriptscriptstyle\triangle}{=}\{1,5,12\}$  means "the set  $\mathcal A$  contains the elements  $1,5,$ and 12".
- Membership  $x \in A$  means "x is an element of A".
- Cardinality  $|A|$  denotes the number of elements in A.

## Example

- 12  $\in$  A: 12 is an element of A.
- 7  $\notin$  A: 7 is not an element of A.

$$
\bullet \ |A|=3.
$$

## Question

$$
\text{Is it true that } \{1,5,12\} = \{5,12,1\} = \{1,1,1,12,5\}?
$$

# Equality and Containment

# Definition

A is a subset of B if every element in A is also an element in  $B$ . Notation:  $A \subseteq B$ .

#### Definition

Two sets A and B are equal if  $A \subseteq B$  and  $B \subseteq A$ . Notation:  $A = B$ .

Definition (strict containment)

$$
A\subsetneq B \Leftrightarrow A\subseteq B \text{ and } A\neq B.
$$

## Example

 $U \triangleq \mathbb{R}$ 

$$
\bullet\; A\stackrel{\triangle}{=}\{1,\pi,4\}
$$

- $\bullet$  B is the interval  $[1, 10]$
- $\bullet$   $A \subset B$ .



The empty set is the set that does not contain any element. It is usually denoted by  $\emptyset$ .

The empty set is a very important set (as important as the number zero).

#### Claim

- $\bullet \forall x \in U : x \notin \emptyset$
- $\bullet \ \forall A \subseteq U : \emptyset \subseteq A$
- $\bullet \ \forall A \subseteq U : A \cup \emptyset = A$
- $\bullet \ \forall A \subseteq U : A \cap \emptyset = \emptyset.$

The power set of a set A is the set of all the subsets of A. The power set of A is denoted by  $P(A)$  or  $2<sup>A</sup>$ .

### Example

The power set of  $A\stackrel{\scriptscriptstyle\triangle}{=} \{1,2,4,8\}$  is the set of all subsets of  $A,$ namely,

$$
P(A) = \{ \emptyset, \{1\}, \{2\}, \{4\}, \{8\}, \\
 \{1,2\}, \{1,4\}, \{1,8\}, \{2,4\}, \{2,8\}, \{4,8\}, \\
 \{1,2,4\}, \{1,2,8\}, \{2,4,8\}, \{1,4,8\}, \\
 \{1,2,4,8\}.
$$

### Question

What is the power set of the empty set  $P(\emptyset)$ ?

#### Question

What is the power set of the power set of the empty set  $P(P(\emptyset))$ ?

# Claim

- $B \in P(A)$  iff  $B \subseteq A$ .
- $\bullet \ \forall A : \emptyset \in P(A)$
- If A has n elements, then  $P(A)$  has  $2^n$  elements. (to be proved)

We can pair elements together to obtain ordered pairs.

# **Definition**

Two objects (possibly equal) with an order (i.e., the first object and the second object) are called an ordered pair.

Notation: The ordered pair  $(a, b)$  means that a is the first object in the pair and  $b$  is the second object in the pair.

# Example

- names of people (first name, family name)
- coordinates of points in the plane  $(x, y)$ .

Equality:  $(a, b) = (a', b')$  if  $a = a'$  and  $b = b'$ .

Coordinates: An ordered pair  $(a, b)$  has two coordinates. The first coordinate equals a, the second coordinate equals b.

The Cartesian product of the sets  $A$  and  $B$  is the set

$$
\mathsf{A}\times\mathsf{B}\stackrel{\scriptscriptstyle\triangle}{=}\{(a,b)\mid a\in\mathsf{A}\;\text{and}\;b\in\mathsf{B}\}.
$$

Every element in a Cartesian product is an ordered pair. We abbreviate  $A^2 \stackrel{\triangle}{=} A \times A$ .

#### Example

Let 
$$
A \stackrel{\triangle}{=} \{0, 1\}
$$
 and  $B \stackrel{\triangle}{=} \{1, 2, 3\}$ . Then,

$$
A \times B = \{(0,1), (0,2), (0,3), (1,1), (1,2), (1,3)\}
$$

### Riddle

Who invented the Cartesian product? (hint: same person invented analytic geometry)

#### Example

The Euclidean plane is the Cartesian product  $\mathbb{R}^2$ . Every point in the plane has an  $x$ -coordinate and a  $y$ -coordinate. Thus, a point  $p$ is a pair  $(p_x, p_y)$ . For example, the point  $p = (1, 5)$  is the point whose x-coordinate equals 1 and whose  *coordinate equals 5.* 

A k-tuple is a set of k objects with an order. This means that a k-tuple has k coordinates numbered  $\{1,\ldots,k\}$ . For each coordinate i, there is object in the ith coordinate.

Alternatively, a  $k$ -tuple is a sequence of  $k$  elements.

- An ordered pair is a 2-tuple.
- $(x_1, \ldots, x_k)$  where  $x_i$  is the element in the *i*th coordinate.
- Equality: compare in each coordinate, thus,  $(x_1, \ldots, x_k) = (x'_1, \ldots, x'_k)$  if and only if  $x_i = x'_i$  for every  $i \in \{1, \ldots, n\}.$

The Cartesian product of the sets  $A_1, A_2, \ldots, A_k$  is the set of all *k*-tuples  $(a_1, \ldots, a_k)$ , where  $a_i \in A_i$ .

$$
A_1 \times A_2 \times \cdots \times A_k \stackrel{\triangle}{=} \{ (a_1, \ldots, a_k) \mid a_i \in A_i \text{ for every } 1 \leq i \leq k \}.
$$

If  $A = A_1 = \cdots = A_k$ , then we abbreviate:

$$
A^k \stackrel{\triangle}{=} A_1 \times A_2 \times \cdots \times A_k
$$

- $\mathbb{R}^3 =$  3-dimensional Euclidean space
- $\mathbb{N}^{12}=$  all sequences of natural numbers that consist of 12 elements.

# De Morgan's Law



Figure: Venn diagram for  $U \setminus (A \cup B) = \overline{A} \cap \overline{B}$ .

# Theorem (De Morgan's Laws)

$$
U \setminus (A \cup B) = \overline{A} \cap \overline{B}
$$
  

$$
U \setminus (A \cap B) = \overline{A} \cup \overline{B}.
$$

To be proved in chapter on Propositional Logic...

A subset  $R \subseteq A \times B$  is called a binary relation.

- Relation of matches between teams in a soccer league. (Liverpool, Chelsea) means that Liverpool hosted the match. Thus the matches (Liverpool,Chelsea) and (Chelsea,Liverpool) are different matches.
- Let  $R \subseteq \mathbb{N} \times \mathbb{N}$  denote the binary relation "smaller than and not equal" over the natural number. That is,  $(a, b) \in R$  if and only if  $a < b$ .

$$
R \triangleq \{(0,1), (0,2), \ldots, (1,2), (1,3), \ldots\}.
$$

A function is a binary relation with an additional property.

## Definition

A binary relation  $R \subseteq A \times B$  is a function if for every  $a \in A$  there exists a unique element  $b \in B$  such that  $(a, b) \in R$ .

A function  $R \subseteq A \times B$  is usually denoted by  $R : A \rightarrow B$ . The set A is called the domain and the set  $B$  is called the range. Lowercase letters are usually used to denote functions, e.g.,  $f : \mathbb{R} \to \mathbb{R}$ denotes a real function  $f(x)$ .

# functions (cont.)

Consider relations  $R_1, R_2, R_3, R_4 \subseteq \{0, 1, 2\} \times \{0, 1, 2\}$ :

$$
\begin{array}{rcl} R_1 & \stackrel{\triangle}{=} & \left\{ (1,1) \right\}, \\ R_2 & \stackrel{\triangle}{=} & \left\{ (0,0), (1,1), (2,2) \right\}, \\ R_3 & \stackrel{\triangle}{=} & \left\{ (0,0), (0,1), (2,2) \right\}, \\ R_4 & \stackrel{\triangle}{=} & \left\{ (0,2), (1,2), (2,2) \right\}. \end{array}
$$

- The relation  $R_1$  is not a function.
- $\bullet$  R<sub>2</sub> is a function.
- The relation  $R_3$  is not a function.
- $\bullet$  The relation  $R_4$  is a constant function.
- $\bullet$  R<sub>2</sub> is the identity function.

- $M \stackrel{\triangle}{=}$  set of mothers.
- $C \triangleq$  set of children.
- $P \stackrel{\triangle}{=} \{ (m, c) \mid m \text{ is the biological mother of } c \}.$
- $Q \triangleq \{ (c, m) | c \text{ is a child of } m \}.$
- $P \subseteq M \times C$  is a relation (usually not a function)
- $\bullet$   $Q \subseteq C \times M$  is a function.

Let  $f : A \rightarrow B$  and  $g : B \rightarrow C$  denote two functions. The composed function  $g \circ f$  is the function  $h : A \rightarrow C$  defined by  $h(a) \stackrel{\triangle}{=} g(f(a)),$  for every  $a \in A$ .

Note that two functions can be composed only if the range of the first function is contained in the domain of the second function.



#### Lemma

Let  $f : A \rightarrow B$  denote a function, and let  $A' \subseteq A$ . The relation  $R\stackrel{\scriptscriptstyle\triangle}{=} (A'\times B)\cap f$  is a function  $R:A'\to B.$ 

 $R$  is called the restriction of  $f$  to the domain  $A'$ .

Let f and g denote two functions. g is an extension of f if  $f \subseteq g$ (every ordered pair in f is also an ordered pair in  $g$ ).

#### Claim

If f :  $A \rightarrow B$  and g is an extension of f, then f is a restriction of g to the domain A.

#### Example

 $f : \mathbb{R} \times \{0\} \to \mathbb{R}$  defined by  $f(x, 0) \stackrel{\triangle}{=} |x|.$ 

• 
$$
g: \mathbb{R} \times \mathbb{R} \to \mathbb{R}
$$
 defined by  $g(x, y) \stackrel{\triangle}{=} \sqrt{x^2 + y^2}$ .

Consider a function  $f : A \times B \rightarrow C$  for finite sets A and B. The multiplication table of f is an  $|A| \times |B|$  table. Entry  $(a, b)$ contains  $f(a, b)$ .

#### Example

The multiplication table of the function  $f: \{0,1,2\}^2 \rightarrow \{0,1,\ldots,4\}$  defined by  $f(a,b) \stackrel{\triangle}{=} a \cdot b.$ 1 | 2 0 0 0 0  $1 \parallel 0 \parallel 1 \parallel 2$  $2 \parallel 0 \mid 2 \mid 4$ 

A bit is an element in the set  $\{0, 1\}$ .

$$
\{0,1\}^n \stackrel{\triangle}{=} \overbrace{\{0,1\} \times \{0,1\} \times \cdots \{0,1\}}^n.
$$

Every element in  $\{0,1\}^n$  is an *n*-tuple  $(b_1,\ldots,b_n)$  of bits.

### Definition

An *n*-bit binary string is an element in the set  $\{0,1\}^n$ .

We often denote a string as a list of bits. For example,  $(0, 1, 0)$  is denoted by 010.

- ${0, 1}^2 = {00, 01, 10, 11}.$
- ${0, 1}^3 = {000, 001, 010, 011, 100, 101, 110, 111}.$

A function  $B: \{0,1\}^n \rightarrow \{0,1\}^k$  is called a Boolean function.

Truth values: "true" is 1 and "false" is 0. Truth table: A list of the ordered pairs  $(x, f(x))$ .

## Example

Truth table of the function NOT :  $\{0, 1\} \rightarrow \{0, 1\}$ :

$$
\begin{array}{c|c}\n \times & \text{NOT}(x) \\
\hline\n0 & 1 \\
1 & 0\n\end{array}
$$

# Important Boolean functions

# **Definition**

\n- \n
$$
\text{AND}(x, y) \triangleq \min\{x, y\}.
$$
\n
\n- \n
$$
\text{OR}(x, y) \triangleq \max\{x, y\}.
$$
\n
\n- \n
$$
\text{XOR}(x, y) \triangleq \begin{cases} \n 1 & \text{if } x \neq y \\ \n 0 & \text{if } x = y \n \end{cases}
$$
\n
\n

Truth tables:

$$
\begin{array}{c|c|c|c|c|c|c|c|c} \hline x & y & \text{AND}(x,y) & & x & y & \text{OR}(x,y) \\ \hline 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 1 & 0 & 0 & 1 & 0 & 1 & 1 & 0 & 1 \\ 0 & 1 & 0 & 0 & 1 & 1 & 0 & 1 & 1 \\ 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 & 0 \\ \hline \end{array}
$$

# Important Boolean functions (cont.)

Truth tables:



Multiplication tables:



# Equivalent definitions

# Claim

•  $NOT(x) = 1 - x$ .

• AND
$$
(x, y) = x \cdot y
$$
.

$$
\bullet \ \mathrm{OR}(x,y)=x+y-(x\cdot y).
$$

• 
$$
\text{XOR}(x, y) = \text{mod}((x + y), 2)
$$

# Multiplication tables:

$$
\begin{array}{c|c|c|c|c|c|c|c|c} \text{AND} & 0 & 1 & \text{OR} & 0 & 1 \\ \hline 0 & 0 & 0 & 0 & 0 & 1 \\ 1 & 0 & 1 & 1 & 1 & 1 \\ \end{array} \hspace{0.2cm} \begin{array}{c|c|c|c|c} \text{OR} & 0 & 1 & \text{XOR} & 0 & 1 \\ \hline 0 & 0 & 1 & 0 & 1 \\ 1 & 1 & 1 & 0 & 0 \\ \end{array}
$$

A function  $f : A \times A \rightarrow A$  is a binary operation.

Usually, a binary operation is denoted by a special symbol (e.g.,  $(+, -, \cdot, \div)$ . Instead of writing  $+(a, b)$ , we write  $a + b$ .

#### Definition

A binary operation  $* : A \times A \rightarrow A$  is commutative if, for every  $a, b \in A$ :

$$
a * b = b * a.
$$

#### Example

 $x + y = y + x$ 

$$
\bullet \ \ x \cdot y = y \cdot x.
$$

$$
\bullet \ \ x - y \neq y - x.
$$

A binary operation  $* : A \times A \rightarrow A$  is commutative if, for every  $a, b \in A$ :

$$
a * b = b * a.
$$

#### Riddle

Why do we care about commutative operations in a logic design course? (hint: Suppose we solder 2 wires to a gate, do we care which wire is soldered to which input?)

A binary operation  $* : A \times A \rightarrow A$  is associative if, for every  $a, b, c \in A$ :

$$
(a * b) * c = a * (b * c).
$$

#### Example

•  $(x + y) + z = x + (y + z)$ 

$$
\bullet (x \cdot y) \cdot z = x \cdot (y \cdot z).
$$

$$
\bullet (x-y)-z\neq x-(y-z).
$$

A binary operation  $* : A \times A \rightarrow A$  is associative if, for every a, b,  $c \in A$ :

$$
(a * b) * c = a * (b * c).
$$

#### Riddle

Why do we care about associative operations in a logic design course? (hint: using 2 gates to compute an operation over 3 bits.) Multiplication of matrices is associative but not commutative:

$$
A = \begin{pmatrix} 1 & 0 \\ 0 & 0 \end{pmatrix}, \qquad \qquad B = \begin{pmatrix} 0 & 1 \\ 0 & 0 \end{pmatrix}.
$$

The products  $A \cdot B$  and  $B \cdot A$  are:

$$
A \cdot B = \begin{pmatrix} 0 & 1 \\ 0 & 0 \end{pmatrix}, \qquad \qquad B \cdot A = \begin{pmatrix} 0 & 0 \\ 0 & 0 \end{pmatrix}.
$$

Since  $A \cdot B \neq B \cdot A$ , multiplication of real matrices is not commutative.

#### **Riddle**

Find a Boolean binary function that commutative but not associative.

#### Question

Given a multiplication table of a binary operator  $f : A \times A \rightarrow A$ , how can we check that  $f$  is commutative? Is there in general a faster way than checking all pairs?

#### Question

Given a multiplication table of a binary operator  $f : A \times A \rightarrow A$ , how can we check that  $f$  is associative? Is there in general a faster way than checking all triples?

#### Question

Prove that both min and max are commutative and associative. What does this imply about  $AND$  and  $OR?$ 

# Associative and Commutative Boolean Functions

# Claim

The Boolean functions OR, AND, XOR are commutative and associative.

### Proof.

Follows from the (algebraic) definitions of the functions.

We can extend the AND and OR functions:

$$
\text{AND}_3(X, Y, Z) \stackrel{\triangle}{=} (X \text{ AND } Y) \text{ AND } Z.
$$

Since the AND function is associative we have

$$
(X \text{ AND } Y)
$$
 AND  $Z = X \text{ AND } (Y \text{ AND } Z)$ .

Thus, we omit parenthesis and write  $X$  AND  $Y$  AND  $Z$ . Same holds for or.